

# Levenberg-Marquardt Method for Mathematical Programs with Linearly Complementarity Constraints

# Cong Zhang<sup>1\*</sup>, Limin Sun<sup>1</sup>, Zhibin Zhu<sup>2</sup>, Minglei Fang<sup>3</sup>

<sup>1</sup>Department of Mathematics and Computer Science, Huarui College, Xinyang Normal University, Xinyang, China
 <sup>2</sup>School of Mathematics & Computational Science, Guilin University of Electronic Technology, Guilin, China
 <sup>3</sup>College of Science, Anhui University of Science and Technology, Huainan, China
 Email: \*<a href="mailto:science">zhcopt@126.com</a>

Received 17 June 2015; accepted 17 August 2015; published 20 August 2015

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#### Abstract

In this paper, a new method for solving a mathematical programming problem with linearly complementarity constraints (MPLCC) is introduced, which applies the Levenberg-Marquardt (L-M) method to solve the B-stationary condition of original problem. Under the MPEC-LICQ, the proposed method is proved convergent to B-stationary point of MPLCC.

## **Keywords**

Mathematical Programs with Linear Complementarity Constraints, MPEC-LICQ, B-Stationarity, Levenberg-Marquardt Method

# **1. Introduction**

The mathematical program with equibrium constraints (MPEC) has extensive application in area engineering design and economic model [1]. It has been an active research topic in recent years. In this paper, we consider the mathematical programming problem with linearly complementarity constraints (MPLCC), which is a special case of the MPEC:

<sup>\*</sup>Corresponding author.

How to cite this paper: Zhang, C., Sun, L.M., Zhu, Z.B. and Fang, M.L. (2015) Levenberg-Marquardt Method for Mathematical Programs with Linearly Complementarity Constraints. *American Journal of Computational Mathematics*, **5**, 239-242. http://dx.doi.org/10.4236/ajcm.2015.53020

$$\min f(x, y)$$

$$s.t. \quad Ax \le b$$

$$w = Nx + My + q$$

$$0 \le w \perp y \ge 0$$

$$(1.1)$$

where  $f: \mathbb{R}^{n+m} \to \mathbb{R}$  is twice continuously differential real-valued function;  $A \in \mathbb{R}^{p \times n}$ ,  $N \in \mathbb{R}^{m \times n}$  and  $M \in \mathbb{R}^{m \times m}$  are given matrices; b and q are given p, m dimensional vectors, respectively.

Complementarity constraints in MPEC are known to be difficult to treat. Research work on the MPEC includes the monograph of Luo *et al.* [1] in which Bouligand stationary condition is introduced that provides a comprehensive study on MPEC. Based on different formulations, there are many algorithms such as Fukushima [2], Zhu [3], Zhang [4] [5], Jiang [6], Tao [7], and Jian [8]. Notice that B-stationary condition is a stronger stationary point. Differing from the approaches mentioned above, we directly introduce L-M technique, without any reformulation or relax form, to solve the B-stationary condition of MPLCC (1.1).

The plan of the paper is as follows: in Section 2, some preliminaries and model we used are presented; in Section 3, the algorithm is proposed.

#### 2. Preliminaries

For reader's convenience, we use following notation throughout this paper:

$$z = (x, y, w), \quad s = (x, y), \quad A^{\mathrm{T}} = (a_{1}^{\mathrm{T}}, a_{2}^{\mathrm{T}}, \dots, a_{p}^{\mathrm{T}}),$$
$$b^{\mathrm{T}} = (b_{1}, b_{2}, \dots, b_{p}), \quad L_{1} = \{1, 2, \dots, p\}, \quad I = \{l \in L_{1} : a_{l}x - b_{l} = 0\},$$
$$L_{2} = \{1, 2, \dots, m\}, \quad I_{y} = \{i \in L_{2} : y_{i} = 0\}, \quad I_{w} = \{i \in L_{2} : w_{i} = 0\}.$$

Let F denote the feasible set of problem (1.1).

Now we give two definitions as follow.

**Definition 2.1.** Let  $z^*$  be a feasible point of MPLCC (1.1), we say that MPEC linear independence constraint qualification is satisfied at  $z^*$  if the gradient vectors

	$\left(A_{S}^{\mathrm{T}}\right)$	$\left( N^{\mathrm{T}} \right)$	$\begin{pmatrix} 0 \end{pmatrix}$		0	
	0	M <sup>T</sup>	$\operatorname{diag}(e_{1i})$		0	
	(0)	$\left( -I \right)$			$\operatorname{liag}(e_{2i})$	))
re	$e_{1i} = \begin{cases} 1, \\ 0 \end{cases}$	$i \in I_y$	, , $e_{2i} =$	<u>∫</u> 1,	$i \in I_w,$	

is linearly independent, where  $e_{1i} = \begin{cases} 1, & i \in I_y, \\ 0, & i \in L_2 \setminus I_y, \end{cases}$ ,  $e_{2i} = \begin{cases} 1, & i \in I_w, \\ 0, & i \in L_2 \setminus I_w. \end{cases}$ 

**Definition 2.2.** Under the MPEC-LICQ, a feasible point z is a B-stationary of problem (1.1) if there exist multiplier vectors  $\lambda \in \mathbb{R}^p$ ,  $\mu \in \mathbb{R}^q$  and  $u, v \in \mathbb{R}^m$  such that

$$\nabla f(z) + \begin{pmatrix} A_{s}^{\mathrm{T}} \\ 0 \\ 0 \end{pmatrix} \lambda + \begin{pmatrix} N^{\mathrm{T}} \\ M^{\mathrm{T}} \\ -I \end{pmatrix} \mu + \begin{pmatrix} 0 \\ \mathrm{diag}(e_{1i}) \\ 0 \end{pmatrix} \mu + \begin{pmatrix} 0 \\ 0 \\ \mathrm{diag}(e_{2i}) \end{pmatrix} v = 0,$$
(2.1)

$$\lambda \ge 0, \quad z \in F, \quad \lambda^{\mathrm{T}} (Ax - b) = 0,$$

$$(2.2)$$

$$u_i = 0, \quad i \in L_2 \setminus I_y, \tag{2.3}$$

$$v_i = 0, \quad i \in L_2 \setminus I_w, \tag{2.4}$$

$$u_i = 0, \quad v_i = 0, \quad i \in I_v \cap I_w.$$
 (2.5)

As we know, most of the works on MPLCC want to get the B-stationary point of problem (1.1), so we also put emphasis on trying to construct a method to obtain the B-stationary of MPLCC (1.1). Now we rewrite the conditions (2.1)-(2.5) in term of lagrange multipliers as follow:

$$Q(\Omega) = \begin{pmatrix} \nabla f(z) + \begin{pmatrix} A_s^{\mathrm{T}} \\ 0 \\ 0 \end{pmatrix} \lambda + \begin{pmatrix} N^{\mathrm{T}} \\ M^{\mathrm{T}} \\ -I \end{pmatrix} \mu + \begin{pmatrix} 0 \\ \operatorname{diag}(e_{1i}) \\ 0 \end{pmatrix} \mu + \begin{pmatrix} 0 \\ 0 \\ \operatorname{diag}(e_{2i}) \end{pmatrix} \nu \\ \operatorname{diag}(e_{2i}) \end{pmatrix} \nu \\ = 0 \qquad (2.6)$$
$$\begin{pmatrix} \lambda_i (a_i x - b_i) \\ Nx + My + q - w \\ u_j y_j \\ v_j w_j \end{pmatrix}$$

subject to:

$$Ax \le b, \quad y \ge 0, \quad w \ge 0, \quad \lambda \ge 0, \quad y_i w_i \le 0, \tag{2.7}$$

and

$$v_l \ge 0$$
 and  $u_l \ge 0$  when  $y_l = w_l = 0$  for some  $l \in L_2$ , (2.8)

where  $\Omega = (z, \lambda, \mu, u, v), \quad j \in L_2$ .

**Remark:** In (2.7) we replace  $y_i w_i = 0$  with  $y_i w_i \le 0$ , because it will be convenient for our computing.

## 3. The Description of Algorithm

Without any reformulation and relaxing techniques, we now use L-M method to solve the nonlinear systems (2.6). Firstly, let J be the Jacobian of  $G(\Omega)$  at  $\Omega$ . For an approximate solution  $z^k$  of (2.6), in order to produce an improving direction, we consider the following system of linear equations

$$\begin{pmatrix} J_k^{\mathrm{T}} J_k + \sigma_k I \end{pmatrix} d = -J_k^{\mathrm{T}} G(\Omega)$$

$$\sigma_k = \theta \| G_k \| + (1 - \theta) \| J_k^{\mathrm{T}} G_k \|,$$

$$(3.1)$$

where  $G_k = G(\Omega^k)$ ,  $\theta$  is a constant.

**Lemma 3.1.** The coefficient matrix of (L - M) is positive definite, and furthermore, (L - M) method has unique solution.

According to the constraint conditions, we now find a step length for current iterated point. First, we consider computing the step length of  $(x, y, w, \lambda)$ . In the first place, for each constraint in (2.7), we should use the  $\Omega^k$  and  $d^k$  to computer a step length:

$$\alpha_{xi} = \begin{cases} 1, & a_i d_x^k \le 0, \\ \min\left(1, \max\left(0, -\frac{a_i x - b_i}{a_i d_x^k}\right)\right), & a_i d_x^k > 0. \end{cases}$$

$$\alpha_x = \min\left(\alpha_{xi}, i \in L_1\right)$$
(3.2)

where  $d_x^k$  is the element of  $d^k$ . Similar to the discussion of step length about *x*, we can obtain the step length  $\alpha_{1y}, \alpha_{1w}, \alpha_{\lambda}$  about  $(y, w, \lambda)$ .

As to calculating the step length for the constraint  $y_j w_j \le 0$ , we get the solution to the equation  $(y_j + \alpha d_y^k)(w_j + \alpha d_w^k) = 0$  with  $\alpha$  as its variable, then  $\alpha_j$  is as follows:

$$\alpha_{j} = \begin{cases} \max(m_{1}, m_{2}), & \text{the equation has two solutions,} \\ \min(1, \max(0, \alpha)), & \text{the equation has one solution and } d_{y}^{k} d_{w}^{k} < 0, \\ 1, & \text{otherwise,} \end{cases}$$
(3.3)

$$\alpha_{2w} = \alpha_{2y} = \min(\alpha_j, j \in L_2),$$

so

$$\alpha_{y} = \min(\alpha_{1y}, \alpha_{2y}), \ \alpha_{w} = \min(\alpha_{1w}, \alpha_{2w})$$

Secondly, we will consider the step length of  $(\mu, u, v)$ . Based on the step length that we obtain above, we can compute the value of  $y_{new}$ ,  $w_{new}$ . If there is some *i* that  $(y_{new})_i = (y_{new})_i = 0$ , then the step length of corresponding variables  $u_i, v_i$  is obtained by the same way in (3.2) in order to satisfy the constraints (2.8); otherwise the step lengths of u, v are set to 1. The step length of  $\mu$  is set to 1. In this paper, we take  $\|G(\Omega)\|^2$  as the merit function.

**Lemma 3.2.** Let d be computed from (3.1), then  $d^{\mathrm{T}}\nabla \|G(\Omega)\|^2 \leq 0$ . Proof. In view of Equation (3.1) and the positive definition of matrix  $(J_k^{\mathrm{T}}J_k + \sigma_k I)$ , we have

$$d^{\mathrm{T}}\nabla \left\|G(\Omega)\right\|^{2} = 2d^{\mathrm{T}}J_{k}G_{k} = -2d^{\mathrm{T}}\left(J_{k}^{\mathrm{T}}J_{k} + \sigma_{k}I\right)d \leq 0.$$

Now we present the algorithm.

## Algorithm A:

Step 0: Given a feasible initial point  $\Omega$ , let k = 1; Step 1: If  $\|G(\Omega)\|^2 < \epsilon$ , then stop; else get the  $d^k$  for (3.1); Step 2: Compute the step length  $\theta_k$ ;

Step 3:  $\Omega^{k+1} = \Omega^k + \operatorname{diag}(\theta_k) d^k$ , go to Step 1, where  $\theta_k = (\alpha_x, \alpha_y, \alpha_w, 1, \alpha_\lambda, \alpha_u, \alpha_y)^T$ .

**Theorem 3.1.** Suppose that  $\Omega$  is generated by Algorithm A and converges to  $\overline{\Omega}$ ; if  $z^k \in F$  for infinitely many k, let the MPEC-LICQ hold on  $\overline{z}$ , then  $\overline{z}$  is a B-stationary point of problem (1.1).

Proof. From the construction of the algorithm, we have  $z^k \in F$  for sufficient large k and  $\overline{z} \in F$ . And because the MPEC-LICQ holds on  $\overline{z}$ , then  $\overline{z}$  is a B-stationary point of problem (1.1).

#### Funding

This work was supported in part by the National Natural Science Foundation (No. 11361018), the Natural Science Foundation of Guangxi Province (No. 2014GXNSFFA118001), Key Program for Science and Technology in Henan Education Institution (No. 15B110008) and Huarui College Science Foundation (No. 2014qn35) of China.

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